Verifiable Binary Lifting

Joe Hendrix, Andrew Kent and Simon Winwood Galois, Inc HCSS 2021



What is Decompilation?

execution.

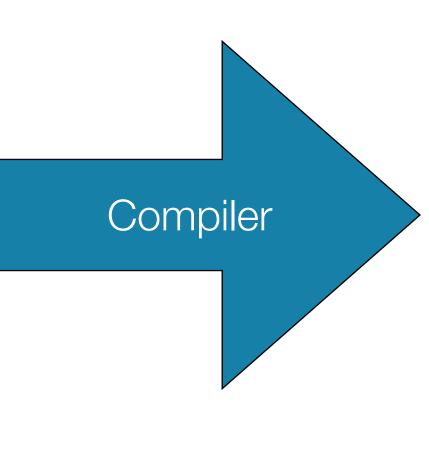
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```
uint64_t fib(uint64_t x) {
    if (x <= 1) {</pre>
           return x;
     } else {
           return fib(x-1)+fib(x-2);
     }
}
```

A decompiler reverses steps in this translation

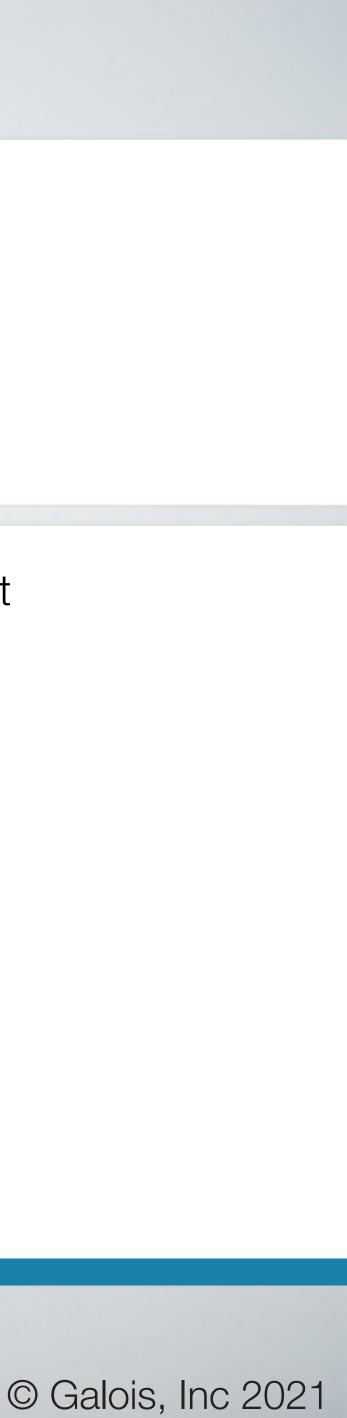
A **compiler** translates code written in a high-level language into a low level language for efficient



0000000000201000 fib:

000000000		
201000:	55	pushq %rbp
201001:	4889e5	movq %rsp,
201004:	4883ec20	subq \$32,
201008:	48897df0	movq %rdi,
20100c:	48837df001	cmpq \$1, -
201011:	0f870d000000	ja 13 <f< th=""></f<>
201017:	488b45f0	movq -16(%
20101b:	488945f8	movq %rax,
20101f:	e934000000	jmp 52 <f< th=""></f<>
201024:	488b45f0	movq -16(%
201028:	482d01000000	subq \$1, %
20102e:	4889c7	movq %rax,
201031:	e8cafffff	callq –54
201036:	488b4df0	movq -16(%
20103a:	4881e902000000	subq \$2, %
201041:	4889cf	movq %rcx,
201044:		movq %rax,
201048:	e8b3ffffff	callq
20104d:	488b4de8	movq -24(%
201051:	4801c1	addq %rax,
201054 :	48894df8	movq %rcx,
201058:	488b45f8	movq −8(%r
20105c:	4883c420	addq \$32,
201060:	5d	popq %rbp
201061:	c3	retq

%rbp %rsp -16(%rbp) -16(%rbp) fib+0x24> %rbp), %rax -8(%rbp) fib+0x58> %rbp), %rax srax %rdi <fib> %rbp), %rcx %rcx %rdi -24(%rbp) -77 <fib> %rbp), %rcx %rcx -8(%rbp) rbp), %rax %rsp



Who uses decompilers?

Decompilers commonly used by reverse engineers to understand a program.

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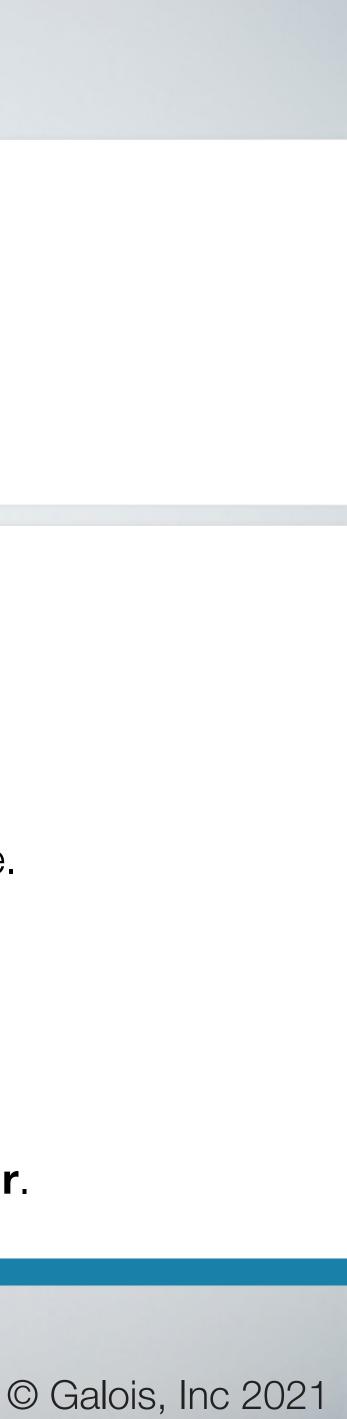
- Decompile into a language understandable by people.
- - structure and local variables.
 - Much more information is lost when compiling with **optimization**.

Engineer works with the decompiler to translate code into idiomatic code.

Without hints or existing source to target, it is generally impossible to recover the original source.

Information lost includes all the structure within function bodies such as original control flow

More recent programs are aimed at using decompilers for program transformation and repair.



Decompilation for Program Transformation

- Researchers are increasingly looking at using decompilers to transform programs.
 - Patch code with vulnerabilities.

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- Extract functionality from legacy code for use in new applications.
- Apply new compiler optimizations or insert security checks into legacy applications.
- Port a program from one platform to another, e.g. x86 to WASM.
- emphasis on programmer understanding.

These new applications place greater emphasis on **program correctness** and may have less

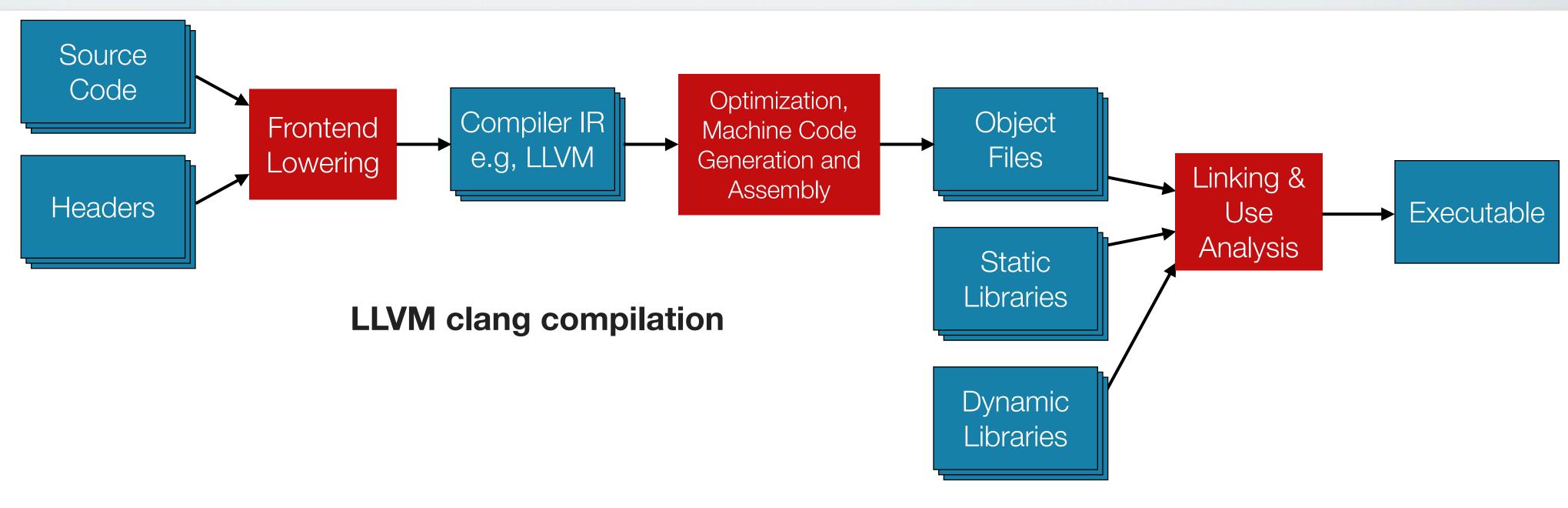


Compilation Toolchain

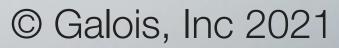
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Decompilation needs to reverse these steps.





Decompilation/Binary Lifting Tools

There are many such tools available:

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- The problem space is very large:
 - Several large and complex instruction set architectures: e.g., x86, x86_64, ARM 8A/7M, PPC, RISCV, ...
 - Variety of operating systems and executable formats: PE, Elf, Macho

Ghidra, IDA/Hex-Rays, Binary Ninja, McSema, RetDec, JEB, Grammatech, Phoenix, reopt

Language and toolchain specific features: e.g., GNU extensions, C++ vtables, eh_frame, ...

Lack of specs for many features, and executables often out of spec (but still "work").



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Related Work Assurance of Decompilation

Scalable Validation of Binary Lifters
 Sandeep Dasgupta, Sushant Dinesh, Deep PLDI, 2020

Evolving Exact Decompilation

Eric Schulte, Jason Ruchti, Matt Noonan, David Ciarletta, Alexey Loginov Workshop on Binary Analysis Research, 2018

Sandeep Dasgupta, Sushant Dinesh, Deepan Venkatesh, Vikram S. Adve, Christopher W. Fletcher



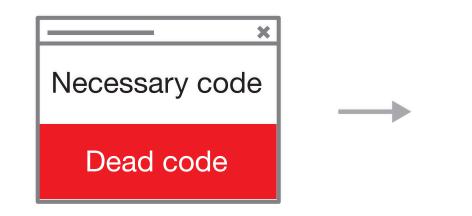
Program Recompilation

• We have implemented an end-to-end recompilation tool called **reopt**.

Application

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Written in a modular fashion so components can be used on other use cases.
 Core binary analysis component also used for verification of machine code.





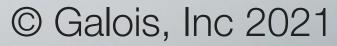
Three Step Process

Three Step Process

- 1. Decompilation
- 2. Recompilation
- 3. Relinking

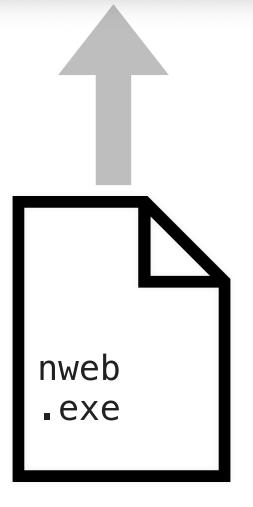
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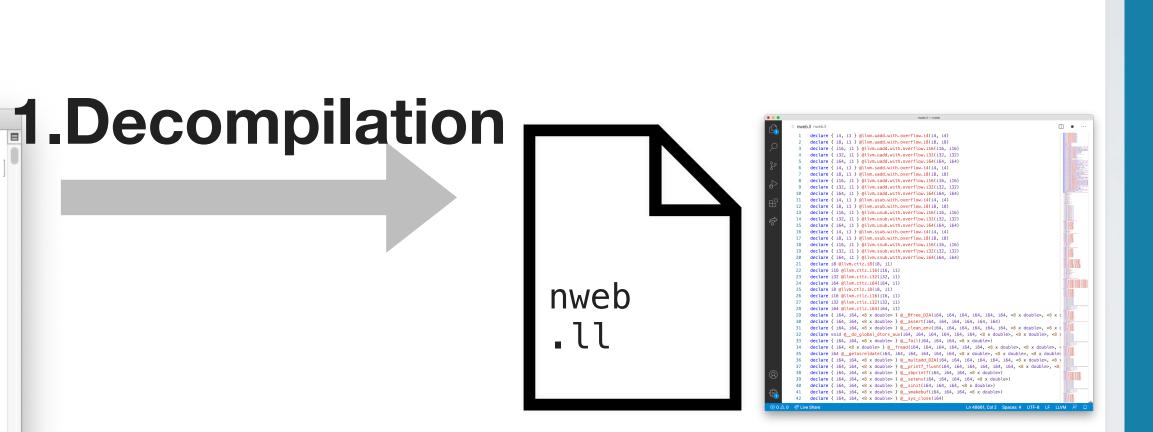
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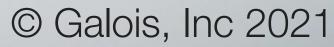




```
reopt — andrew@000385-andrew — ..os/VADD/reopt — -zsh — 80×16
Last login: Fri Oct 2 13:01:43 on ttys001
→ reopt git:(master) × cabal run reopt -- nweb23_static_freebsd
Up to date
Analyzing function: 0x400138 (_init)
Analyzing function: 0x400150 (_start)
Analyzing function: 0x4001f0 (__do_global_dtors_aux)
Analyzing function: 0x400240 (frame_dummy)
Analyzing function: 0x400290 (logger)
Analyzing function: 0x400480 (web)
Analyzing function: 0x400830 (main)
Analyzing function: 0x400c40 (_
                                _bswap16_var)
Analyzing function: 0x400c60 (__tls_get_addr)
Analyzing function: 0x400c70 (_init_tls)
Analyzing function: 0x400d80 (_rtld_allocate_tls)
Analyzing function: 0x400e60 (_rtld_free_tls)
Analyzing function: 0x400e90 (sleep)
```





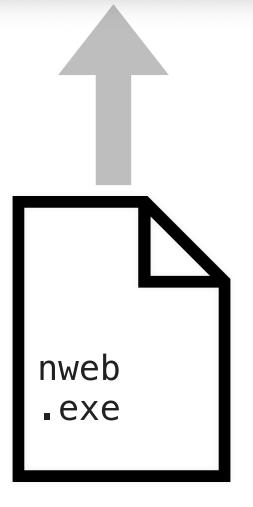


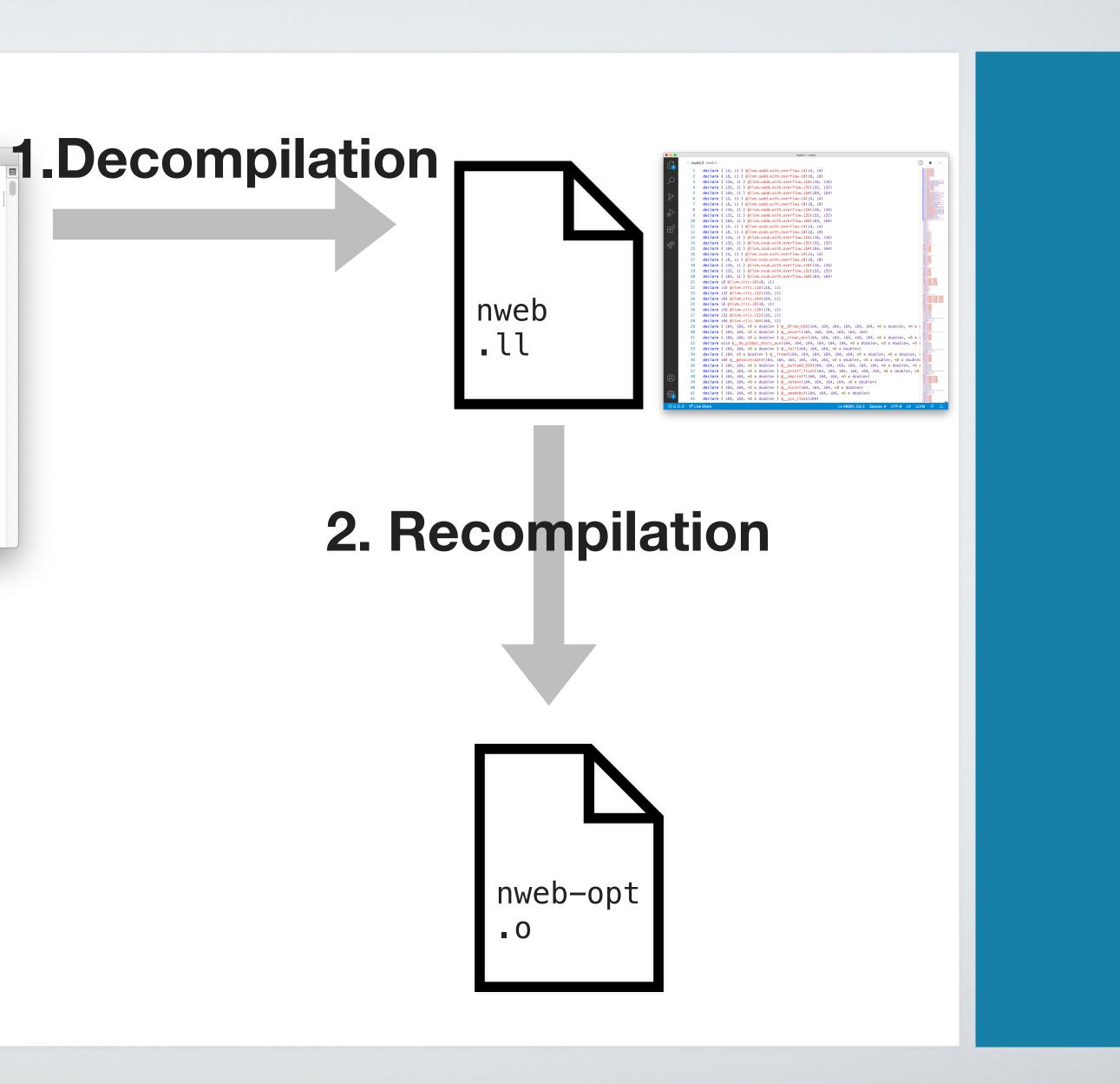


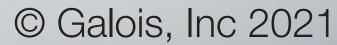
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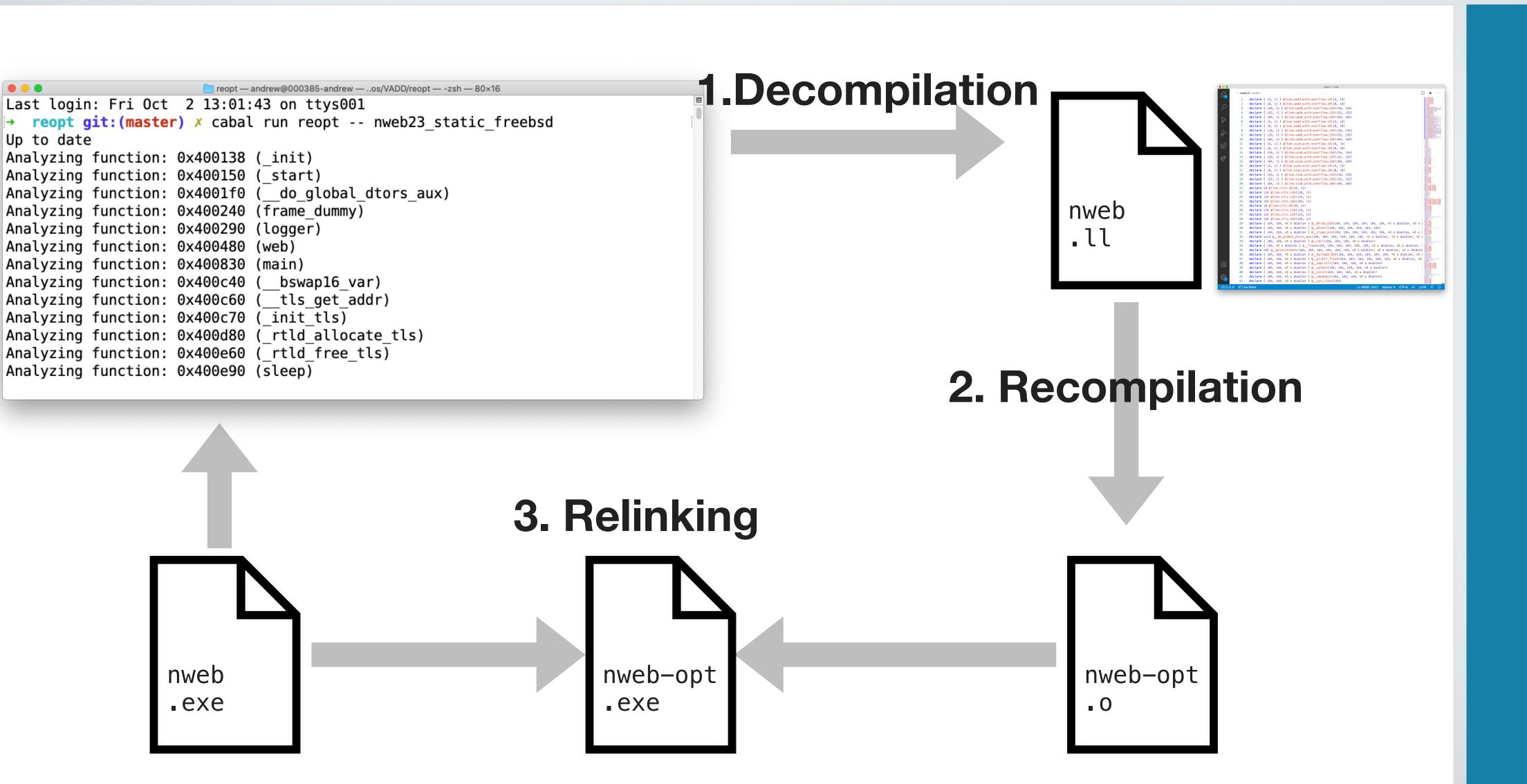


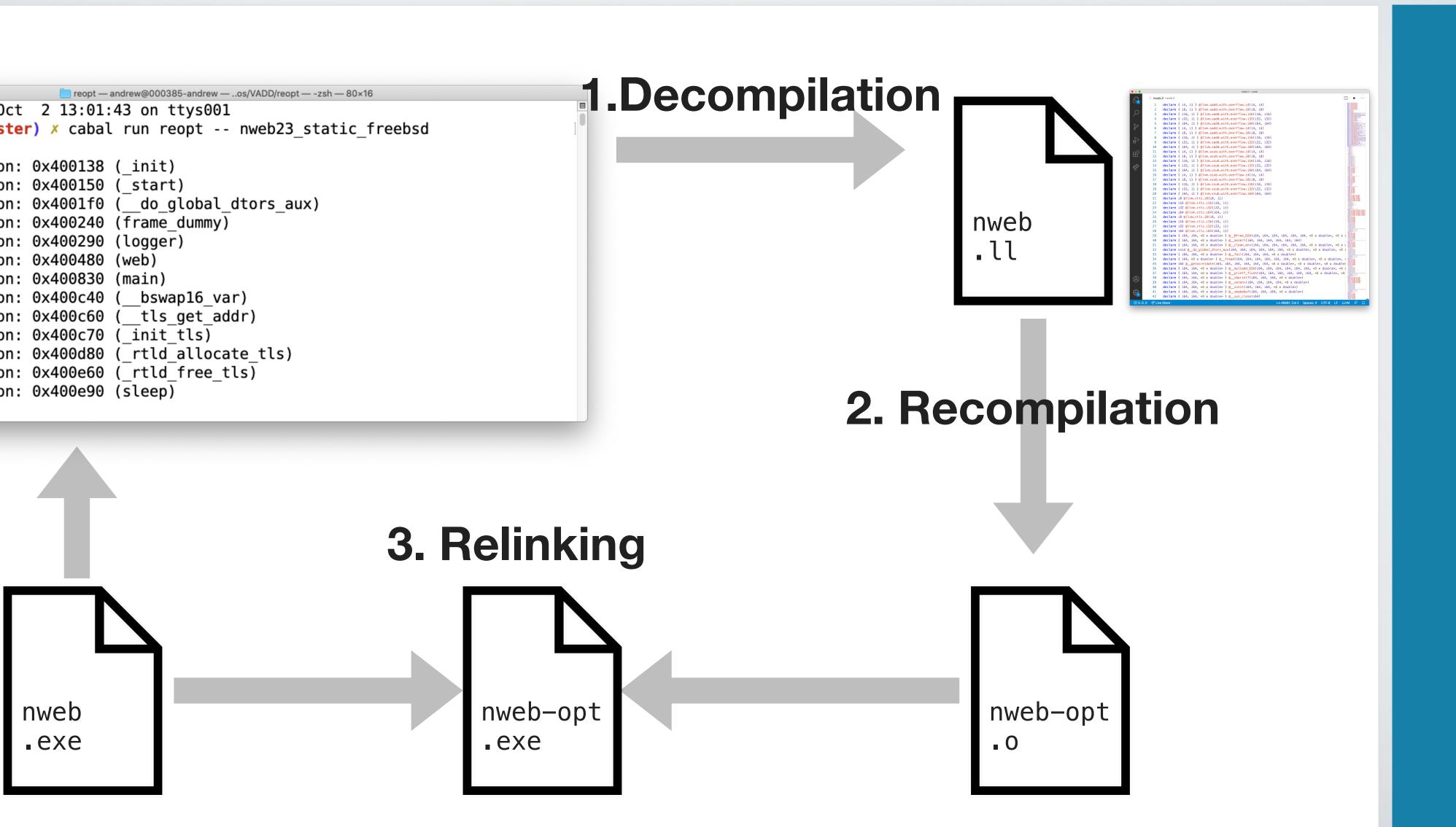












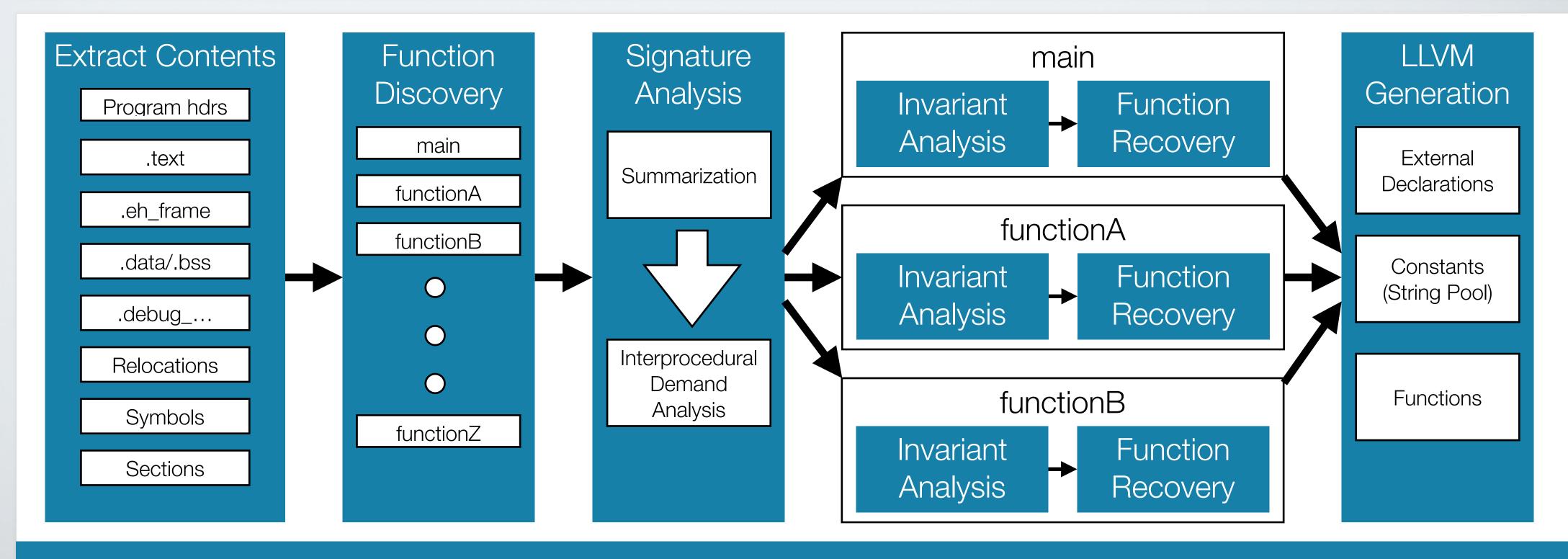






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Decompilation Pipeline



Supported Features

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- Analysis component has large coverage of x86_64, ARM and PowerPC ISAs
 - x86_64 includes significant SSE/AVX support, some FPU/MMX support.
- Invariant inference and LLVM generation limited to x86_64 ISA.
 - x86_64 coverage is a more limited.
- Static and dynamically linking executables.
- Extracting information from debug and .eh_frame data.



15 Handling Failures

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Pipeline steps may fail, and reopt recovers as much as it can.

- I. Elf files sometimes inconsistent or use extensions we do not support.
- 2. Discovery may fail to find all functions or spurious code.
 - Funcitions only indirectly referenced in structs primary failure to find.
 - Spurious code from no-return functions.
- 3. Function argument analysis is main limitation currently.
 - Lack mechanism for inferring arguments to externally linked functions.
- Invariant analysis and function recovery do not support full instruction set. 4.



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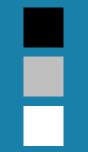
Interoperability

- Can export intermediate results at each s pipeline.
- VSCode extension under development.
- Binary analysis plugins for Ghidra and Bir developed on previous projects.

						[Extension Development Host] - test_add_diet_IId.dis — examp
	Q	REOPT: PROJECT	≣ test_ado	d_diet_lld.dis 9+, U ×		
	لك	Create reopt project	≣ test_a	dd_diet_lld.dis		
	Q	Open reopt project	29	20136c: 48 8d 54 fe 08	lea	rdx,[rsi+rdi*8+0x8]
	~		30		mov	QWORD PTR [rip0x1418],rdx
	90	Close reopt project	31	201378: e8 95 00 00 00	call	pc+95
	908	test_add_diet_lld	32	20137d: 48 89 c7	mov	rdi,rax
			33 34	201380: e8 01 00 00 00 201385: f4	call hlt	pc+1 Could not determine arguments for call to
otogo of	æ	Generate disassembly	34	201386: b0 3c	mov	al,0x3c
stage of		Generate functions	36	201388: b4 00	mov	ah,0x0
	Ē	Generate CFG	37	20138a: Of b7 c0	movzx	eax,ax
	_	Generate LLVM	38	20138d: 49 89 ca	mov	r10, rcx
	B		39	201390: Of 05	syscal	
			40	201392: 48 3d 7c ff ff ff	cmp	rax,-0x84
	.		41	201398: 76 Of	jbe	<pre>pc+f Call targets must be direct calls.</pre>
			42	20139a: f7 d8	neg	eax
	$\langle \mathcal{O} \rangle$		43	20139c: 50	push	rax
			44	20139d: e8 08 00 00 00	call	pc+8
	E		45	2013a2: 59	рор	rcx
inon / Nlinio			46	2013a3: 89 08	mov	DWORD PTR [rax],ecx
inary Ninja	Json		47	2013a5: 48 83 c8 ff	or	rax,-0x1
, ,			48	2013a9: c3	ret	
			49 50	2013aa: 48 c7 c0 fc ff ff ff		rax,-0x4
			50 51	2013b1: 64 48 03 04 25 00 00 2013b8: 00 00	add	rax,QWORD PTR fs:0x0 LLVM backend does not yet
			51	2013ba: c3	ret	
			53	2013bb: 48 8b 07	mov	rax,QWORD PTR [rdi]
			54	2013be: 48 85 c0	test	rax,rax
			55	2013c1: 74 10	je	pc+10
			56	2013c3: 48 39 f0	cmp	rax, rsi
			67	2012-C. 75 05	ina	
			PROBLEMS	0UTPUT DEBUG CONSOLE	TERMINA	L _o
			at ChildProcess.exithandler (child_process.js:312:12)			
			at ChildProcess.emit (events.js:315:20)			
			at maybeClose (internal/child_process.js:1021:16)			
			at Socket. <anonymous> (internal/child_process.js:443:11)</anonymous>			
				ocket.emit (events.js:315:20)		
				ipe. <anonymous> (net.js:674:12</anonymous>		
			at Pipe.callbackTrampoline (internal/async_hooks.js:120:14)			
				owing command errored, scroll		
	£633		/users/v	al/galois/reopt/result-reopt/b	in/reop	tllvmoutput= <u>/Users/val/galois/reopt/vscode-plu</u>

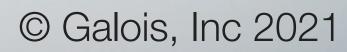
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Verification







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Verification Properties

Recompilation Soundness

Verification Soundness

Every observable execution in the LLVM should be possible in the machine code program.

 $t \in traces(P_{LLVM}) \Rightarrow \exists t' \in traces(P_{MC}), t \equiv t'$

If a property is true of the raised program, then it should be true of the machine code program.



Observational Equivalence

- Our current notion of equivalence is based on event traces.
- Required events include:

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- Writes to non-stack addresses.
- Other operations that may raise signals (e.g., divide-by-zero).
- System calls
- Internally, we make additional equivalence checks for compositional purposes.



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Verification Approaches

1. Build a **verified decompiler** using interactive theorem proving.

2. Use **automation** to check program equivalence after generation.



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Verification Approaches

1. Build a **verified decompiler** using interactive theorem proving.



- Decompilation is an open-ended problem.
- Very complex to implement, and needs continued improvement.

- 2. Use **automation** to check program equivalence after generation.
- Program equivalence is ordinarily undecidable...
- However, the decompiler output is structurally similar to input binary.
- We have developed a compositional approach that checks equivalence of basic blocks using SMT solving.



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Verification Approach

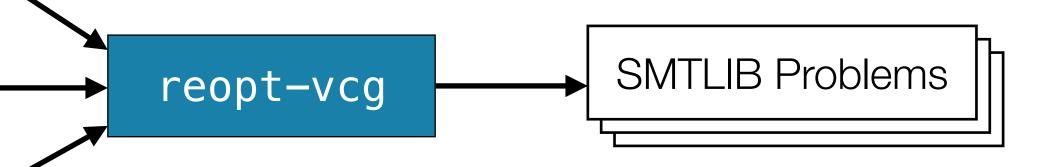
We have implemented a verifier based on translation validation.

Original Binary

Generated Annotations

Generated LLVM

Correctness claim: If all SMTLIB SAT problems are unsat, then the generated LLVM refines the original binary





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VCG Implementation

The implementation of reopt-vcg is independent of reopt itself.

Reopt

- Written in Haskell
- Disassembler based on udis86.
- Custom x86_64 semantics.
- Haskell LLVM-pretty library used for generating LLVM.

reopt-vcg

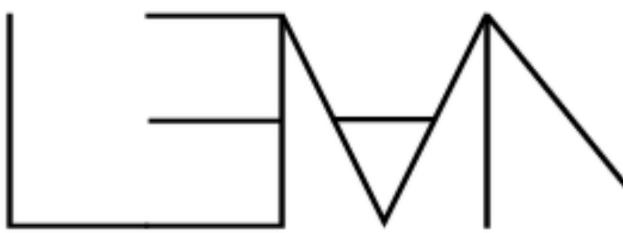
- Written in Lean 4
 - Long-term goal is to verify reopt-vcg.
- Disassembler from LLVMMC.
- UIUC K x86_64 semantics.
- libllvm used for parsing

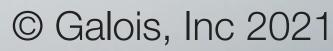


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- Lean 4 is a new programming language and theorem prover.
 - Publicly available but not officially released yet.
 - Functional language based on dependent type theory.
 - Compiles to efficient C code; No cyclic data-structures.
 - Self hosting Largely written in Lean itself.
- Builds on Lean 3
 - Mathlib a large community effort to formalize mathematics.
 - ~450kloc of definitions and theorems as of January 2021.









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Satisfiability Modulo Theories (SMT) Solvers

SMT Solvers can automatically prove theorems involving specific **theories**.

- A theory has one or more types (called sorts in SMT) along with operations.
- There are combination methods that allow theory solvers to work together.
- reopt-vcg uses bitvectors, arrays (with a partial equivalence extension), and uninterpreted functions.
- SMT solving is NP-hard (some theories are more difficult), but a lot of work has gone into making it tractable.

SMT-COMP 2021



https://smt-comp.github.io/



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Compositional Proofs

- smaller proofs.
- Instead of asking:
- We instead ask solvers to answer many questions of the form:
- For a compositional strategy, we need
 - All the assumptions needed to make the statement true.
 - Check that the assumptions hold when jumping from one block to another.

The key to making automation tractable is to decompose the overall equivalence of programs into many

Is LLVM Program P equivalent to machine code program Q?

Is this effect in a LLVM basic block B equivalent to this effect in the machine code?



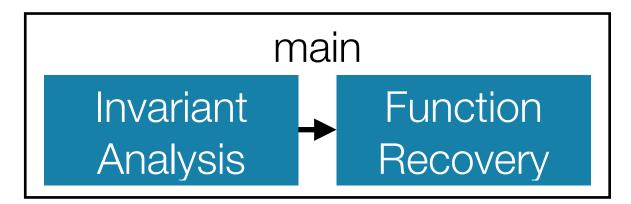
Compositional Properties

Reopt-VCG's compositional strategy enforces

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- Functions respect the ABI (how arguments are passed, callee-saved registers, etc)
- The size of each stack frame is bounded to at most a page and all stack accesses are in bounds.
 - Needed to avoid accessing heap memory via stack pointers.
- Callee saved information is saved and persisted and not modified by the function.





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Getting Reopt

- reopt and reopt-vcg are publicly available under open source libraries.
 - https://github.com/GaloisInc/reopt

You can try it out online through Gitpod, download a Docker image, or use prebuilt binaries.



Some Observations

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- The binary analysis ecosystem has a rich variety of tools.
 - Made much easier by extensive documentation and open-source libraries.
 - Prefer libraries that provide semantics as data or a DSL rather than code.
- Large and complex instruction sets represent a significant but understood challenge.
- Operating system, debug and linker extensions challenging:
 - Lack of consistent documentation
 - Large variety of extensions
 - Implementations make different choices and change regularly.



Thank you

Acknowledgements:

- Guannan Wei of Purdue implemented an early prototype of reopt-vcg.
- Rob Dockins (Galois) wrote the Lean LLVM bindings.
- Leonardo de Moura and Sebastian Ulrich for help with Lean 4.
- The CVC4 team for the partial array equality extension to CVC4.
- Andrei Stefanescu (Galois) and the UIUC K Semantics group for their x86_64 semantics.

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